

# Reflection and its use

## from science to meditation

Review

# Music

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Pierluigi da Palestrina (1525?-1594)

*Ave Verum*

Traditional Buddhist Music (Shakuhachi, Japan)

Igor Stravinsky

*Introitus* (1965) In memoriam for T.S. Eliot

# Contents

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Logic

The Turing Test

Sharp reflection

Reverse reflection

# Propositional Logic

Introduction Rules	Elimination Rules
$\frac{\Gamma, A \vdash B}{\Gamma \vdash (A \rightarrow B)}$	$\frac{\Gamma \vdash (A \rightarrow B) \quad \Gamma \vdash A}{\Gamma \vdash B}$
$\frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash (A \ \& \ B)}$	$\frac{\Gamma \vdash (A \ \& \ B) \quad \Gamma \vdash (A \ \& \ B)}{\frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash (A \vee B) \quad \Gamma, A \vdash C \quad \Gamma, B \vdash C}}$
$\frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash (A \vee B)}$	$\frac{\Gamma \vdash (A \vee B) \quad \Gamma, A \vdash C \quad \Gamma, B \vdash C}{C}$
Start Rule	Absurdum Rule
$\frac{A \in \Gamma}{\Gamma \vdash A}$	$\frac{\Gamma \vdash \perp}{\Gamma, \neg A \vdash \perp}$
	$\neg A := (A \rightarrow \perp)$
	$\frac{\Gamma, \neg A \vdash \perp}{\Gamma \vdash A}$

## Example

$$A \rightarrow (A \rightarrow B) \vdash A \rightarrow B$$

# Predicate Logic

This is all of logic (Aristotle, Boole, Peirce, Frege, Brouwer, Heyting, Gentzen)

	Introduction Rules		Elimination Rules	
$\rightarrow$	$\Gamma, A \vdash B$	$\frac{}{\Gamma \vdash (A \rightarrow B)}$	$\Gamma \vdash (A \rightarrow B)$	$\Gamma \vdash A$
	$\Gamma \vdash A \quad \Gamma \vdash B$	$\frac{}{\Gamma \vdash (A \& B)}$	$\Gamma \vdash (A \& B)$	$\Gamma \vdash (A \& B)$
	$\Gamma \vdash A \quad \Gamma \vdash B$	$\frac{}{\Gamma \vdash (A \vee B)}$	$\Gamma \vdash A$	$\Gamma \vdash B$
$\vee$	$\Gamma \vdash A$	$\Gamma \vdash B$	$\Gamma \vdash (A \vee B)$	$\Gamma, A \vdash C \quad \Gamma, B \vdash C$
	$\frac{}{\Gamma \vdash (A \vee B)}$	$\frac{}{\Gamma \vdash (A \vee B)}$		$C$
$\forall$	$\Gamma \vdash A$	$x \notin \Gamma$	$\Gamma \vdash \forall x.A$	$t$ is free in $A$
	$\Gamma \vdash \forall x.A$		$\Gamma \vdash A[x := t]$	
$\exists$	$\Gamma \vdash A[x := t]$		$\Gamma \vdash \exists x.A \quad \Gamma, A \vdash B$	$x \notin B$
	$\frac{}{\Gamma \vdash \exists x.A}$		$\Gamma \vdash B$	
	Start Rule	Absurdum Rule	Classical Negation	
	$A \in \Gamma$	$\Gamma \vdash \perp$	$\Gamma, \neg A \vdash \perp$	
	$\frac{}{\Gamma \vdash A}$	$\frac{}{\Gamma \vdash A}$	$\neg A := (A \rightarrow \perp)$	
			$\Gamma \vdash A$	

## Example

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$$A \rightarrow (A \rightarrow B) \vdash A \rightarrow B$$

Can be applied in connection with

$$\forall x, y. (x > y) \rightarrow \neg(y > x) \quad \text{axiom}$$

$$\forall x, y. x > y \rightarrow (y > x \rightarrow \perp)$$

$$x > x \rightarrow (x > x \rightarrow \perp)$$

$$x > x \rightarrow \perp$$

$$\neg(x > x)$$

$$\forall x. \neg(x > x) \quad \text{consequence}$$

# The Turing Test

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“Does a machine have consciousness?”

In order to make this question more precise

Turing modified the *simulation test*:

find out while chatting whether two persons are male or female,  
while one is simulating to be the other gender

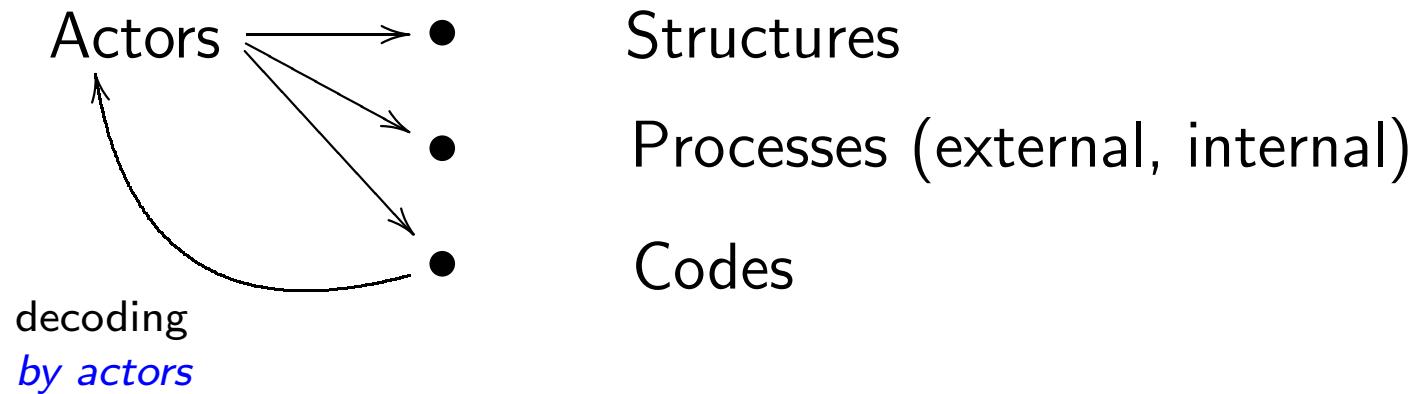
Turing's version: decide between a human and a robot  
if we cannot distinguish correctly, we may say that  
the robot has consciousness

# Sharp reflection

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Reflection

Global feedback



Ordinarily it does not always happen that actors modify the code

In *sharp* reflection it does

# Sharp reflection in language

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## Rebus (in Dutch)



→ stoomboot

b=st

m=t

The following correction process can be captured into words

“Hij bedoeld iets” →

“Hij bedoeld~~e~~ iets” X t →

“Hij bedoelt iets”

# Sharp reflection in Computing

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A *compiler* takes a program  
and transforms it into another one  
directly interpretable by the electronics on a chip

# Reverse reflection

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actor  $\longmapsto$  code

## Reverse reflection in mathematics

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Continuous functions (can write graph without lifting pen from paper)

$$\sin(x)$$

$$\sin(x) + \sin(2x)$$

$$\sin(x) + \sin(2x) + \sin(4x)$$

...

Theorem

$$\forall e \in L. \llbracket e \rrbracket \text{ is continuous in } x$$

Here  $L$  is the language consisting of the names of such functions:

$$\{ \sin, \sin_2, \sin_3, \dots \}$$

closed under arbitrary sums

In order to show that for example  $\sin(x) + \sin(4x)$  is continuous we write it as

$$\sin(x) + \sin(4x) = \llbracket \sin + \sin_4 \rrbracket$$

and apply the theorem. We need to go from the actor to its code!

## The quote mode

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To go from an actor to its code is possible, if we have that code already in one way or another. We may suppose that every mathematical object (actor) is given as its code and is interpreted whenever convenient or necessary.

This actually happens in a version of the programming language LISP and is called the quote-mode of this language.

Increased states of mindfulness also work like this.