

## Week 2

### In-class problems

1. (a) Let  $g, h$  be  $\lambda$ -defined by  $G, H$ . Find a  $F \in \Lambda^\emptyset$  that is  $\lambda$ -defining

$$\begin{aligned} f(0, y) &= g(y) \\ f(x + 1, y) &= h(f(x, y), x, y). \end{aligned}$$

- (b) Lambda define the predecessor (having on 0 the value 0 itself).
2. (a) Show that  $\mathbf{times} \ \mathbf{c}_n \ \mathbf{c}_m =_\beta \mathbf{c}_{n \cdot m}$ , for all  $n, m \in \mathbb{N}$ .  
 (b) Show that  $\mathbf{c}_n \mathbf{c}_m =_\beta \mathbf{c}_{m^n}$ , for all  $n, m \in \mathbb{N}$ .
3. (a) Define the characteristic function  $K_\leq$  of the relation  $\leq$ . Show that it is  $\lambda$ -definable.  
 (b) Show that  $L(n, m) = \lceil {}^n \log m \rceil$  is  $\lambda$ -definable. Here, for  $r \in \mathbb{R}$  the integer  $\lceil r \rceil$  is its ceiling, the least element  $n \in \mathbb{Z}$  such that  $r \leq n$ . [Hint. Use 2(b) and (a).]
4. Give an inductive definition of trees with natural numbers at the leaves.
  - (a) Specify the function mirror on trees. Show that the function  $\mathbf{mirror} \in \Lambda^\emptyset$  (given on the slides) defines it.
  - (b) Specify the function that squares all leaves in a tree, leaving the tree structure the same. Construct a  $\lambda$ -defining term for this function.

5. (Klop) Define

$$\begin{aligned} \$ &\triangleq \lambda abcdefghijklmnopqrstuvwxyzr.r(\text{this is a fixed point combinator}) \\ \mathbf{\text{€}} &\triangleq \$. \end{aligned}$$

Show that  $\mathbf{\text{€}}$  is a reducing fixed point combinator:  $\mathbf{\text{€}}f \rightarrow_\beta f(\mathbf{\text{€}}f)$ .

### Take-home problems

1. Show that there are no  $F_1, F_2 \in \Lambda^\emptyset$  such that

$$F_1(xy) = x \ \& \ F_2(xy) = y.$$

[Hint.  $F_1$  alone doesn't even exist.]

2. (Petrossi) Show that  $\mathbf{C}$  has no nf, where
- $$\begin{aligned} \mathbf{A} &\triangleq \text{SSS} \\ \mathbf{B} &\triangleq \text{SAA} \\ \mathbf{C} &\triangleq \text{BB}. \end{aligned}$$

3. Find a different representation  $n \mapsto \mathbf{d}_n$  of the numbers with  $\mathbf{d}_n \in \Lambda^\emptyset$  such that  $|\mathbf{d}_n| = O(\lg n)$  and there are  $F, G \in \Lambda^\emptyset$  satisfying

$$F\mathbf{d}_n =_\beta \mathbf{c}_n \ \& \ G\mathbf{c}_n =_\beta \mathbf{d}_n.$$

Show that the terms do what is intended.

4. (a) Show that commutativity or associativity in  $\lambda$  does not hold in the following strong sense. Adding to  $\lambda$  one of the following axioms makes it inconsistent.

$$xy = yx \tag{1}$$

$$(xy)z = x(yz) \tag{2}$$

- (b) Adding either (1) or (2) to CL does not make it inconsistent; show that for this one *needs* an axiom scheme like

$$XY = YX \text{ for all terms } X, Y.$$

5. (Barendregt, Klop, Dezani)

- (a) Show that there are  $A, B, C, E \in \Lambda^\emptyset$  that behave like Klein's four-group with multiplication table below. [Hint. First solve (b).]

	$A$	$B$	$C$	$E$
$A$	$E$	$C$	$B$	$A$
$B$	$C$	$E$	$A$	$B$
$C$	$B$	$A$	$E$	$C$
$E$	$A$	$B$	$C$	$E$

- (b) Let  $f: \mathbb{N}^2 \rightarrow \mathbb{N}$  be computable. Construct  $=_\beta$ -distinct closed  $\lambda$ -terms  $X_0, X_1, X_2 \dots$  such that for all  $n, m \in \mathbb{N}$  one has

$$X_n X_m =_\beta X_{f(n,m)}.$$

[Hint. Try  $X_n \triangleq \langle A, \mathbf{c}_n \rangle$ .]